Tuning Trombones, Smoothing Speedbumps:

Why Internet Speeds in Africa are Slow, and What We Can Do About It

Nick Feamster Princeton University

Broadband Connectivity in Africa

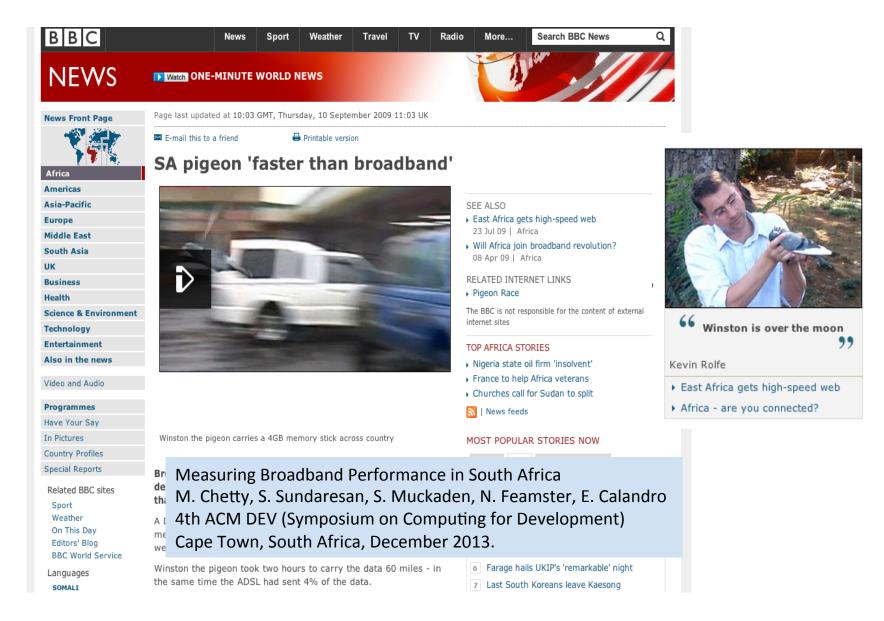


According to ITU in 2013

- 93 million broadband subscriptions
- 27% growth in past 4 years (Highest)
- Broadband associated with economic growth + development

Yet, little is known about performance in Africa and what causes poor performance when it does arise.

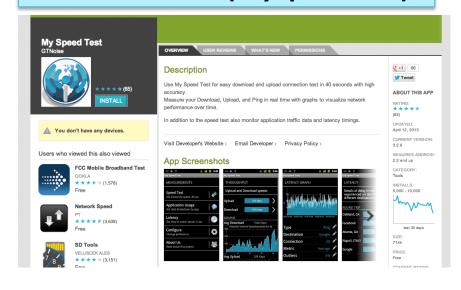
How Well Does Broadband Perform?



From Where To Measure?

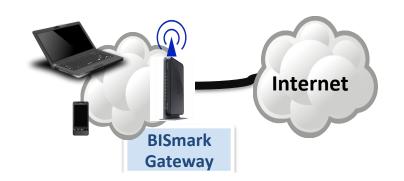


Mobile Handset (MySpeedTest)



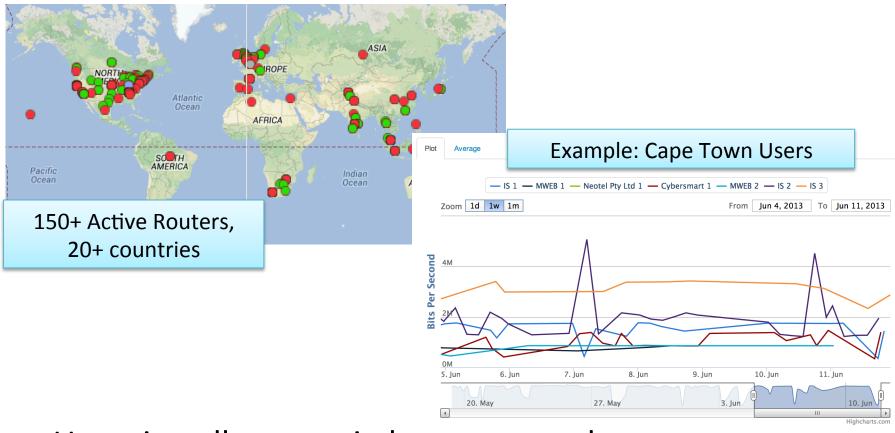
Home Gateway (BISmark)





- Fixed-line measurements
- Dongle-based measurements

BISmark: Measurements from Fixed Line Locations



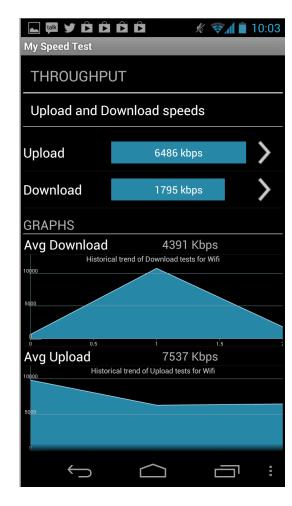
- Users install routers in home networks
- Custom firmware performs periodic measurements
- Can aggregate by country, city, ISP

My Speed Test: Measurements from Mobile Handset

- Periodic latency measurements
- Upload and download throughput tests on demand
- Available in Google Play Store:

http://goo.gl/28tx3





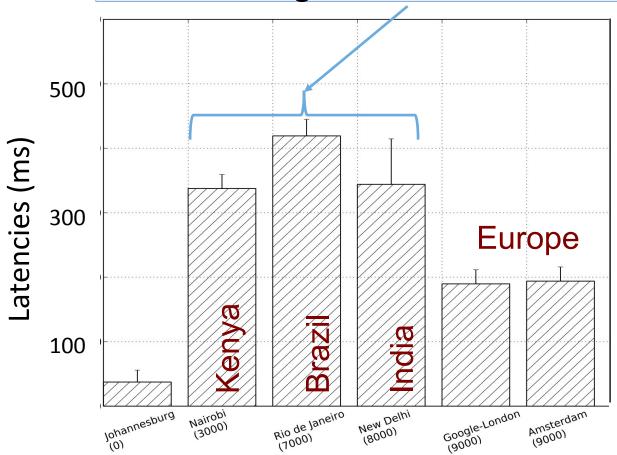




- Performance consistently underperforms with respect to advertised rates
- Mobile broadband consistently achieves higher throughput than fixed broadband
- (Bad) peering can introduce significant latency, introduce fragility in times of failure (e.g., fiber cut)

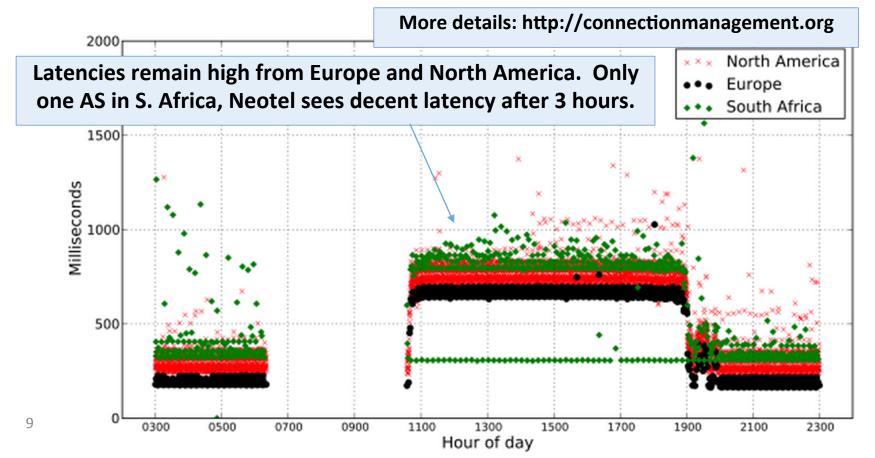
Latencies to Nearby Locations are High

Latencies from South Africa to Kenya, Brazil, India are 2x higher than latencies to Europe.

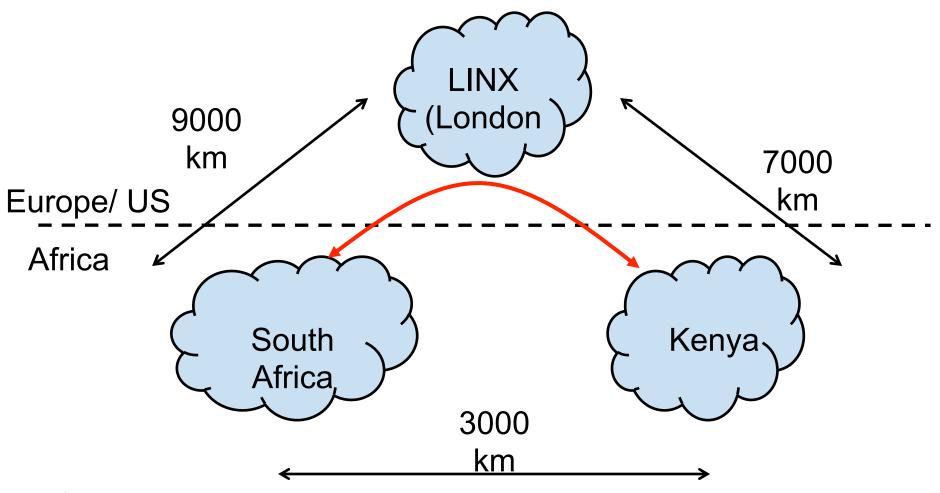


Latencies are Even Higher During Failures

- •March 27, 2013 0620 UTC: SWM4 Fiber Cut
- •All BISmark hosts could not reach KENet for 3+ hours
- •Latencies remain high for another 8+ hours (except for Neotel, in South Africa)



Causes of High Latency: Circuitous Routing Paths



Two Questions

 What is the nature of Internet interconnectivity (between ISPs) in Africa?

 What can be done to reduce latency to common Internet services?

Peering at the Internet's Frontier: A First Look at ISP Interconnectivity in Africa A. Gupta, M. Calder, N. Feamster, M. Chetty, E. Calandro, E. Katz-Bassett Passive and Active Measurement Conference Los Angeles, CA, March 2014.

Two Questions

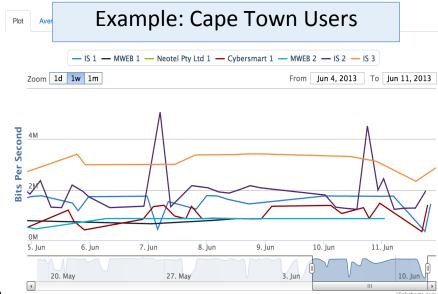
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BISmark: Measurements from Fixed Locations



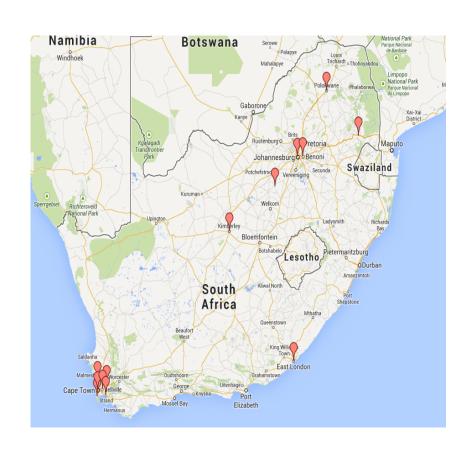


- Users install routers in home networks
- Custom firmware performs periodic measurements
- Can aggregate by country, city, ISP



BISmark Deployment in South Africa

- Periodic latency and throughput measurements
- Traceroutes to explain the cause of path performance
- Router-based deployment
 - 17 home networks, 7 ISPs, all 9 provinces



Destinations for Traceroute Probes

Global M-Lab Servers

Canada Poland Germany Ulraine Kazakhstan Monopolia China South Korea Alperia Libya Fakistan Alperia Libya Fakistan Pakistan Anghanistan Ocean Alperia Libya Formany Alphanistan Pakistan Anghanistan Anghanistan Pakistan Anghanistan Anghanistan Pakistan Anghanistan Anghanistan Pakistan Anghanistan Pakistan Anghanistan Anghanistan Anghanistan Pakistan Anghanistan Anghanistan Anghanistan Anghanistan Pakistan Anghanistan Anghanis

Google Caches in Africa



High Latencies to Nearby Locations...

Cape Town (SA) to M-Lab Johannesburg (SA)

Cape Town (SA) to M-Lab Nairobi (KE)

•••

7, 196.44.0.74, 7.793, South Africa, AS16637

8, 196.223.22.24, 8.338, South Africa, Cape Town IXP

9, 41.164.0.243, 34.679, South Africa, AS36937

•••

14, 196.24.45.146, **92.511**, South Africa, AS2018

8, 209.212.111.201, 199.446, South Africa, AS16637

9, 195.66.225.31, 217.301, United Kingdom, London IXP (LINX)

10, 196.32.209.77, 201.569, South Africa, AS36944

•••

14, 197.136.0.108, **368.107**, Kenya, AS36914

High Latency

... Circuitous Routing Paths

Cape Town (SA) to M-Lab Johannesburg (SA)

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Traffic leaving Africa

Poor ISP Interconnectivity in Africa

- Reasons
 - Local ISPs not present at regional IXPs
 - IXP participants don't peer with each other

- Consequences
 - Local traffic does not stay local
 - Paths leave continent

Local ISPs not Present at Regional IXPs

• ISBs prioritize connecting to European

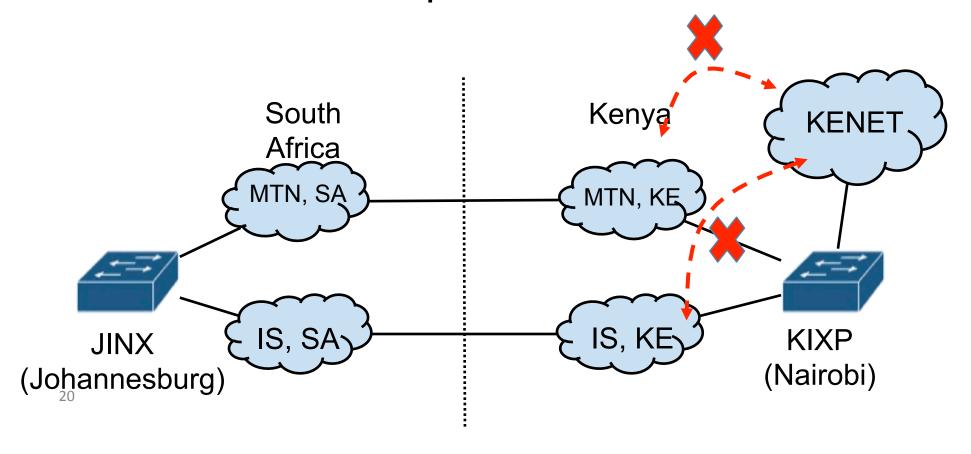
• Fewer incentives to connect at regional ones

South Kenya
Africa

Liquid
Telecom
KIXP
(Nairobi)

Missing Peering Links at Regional IXPs

- Most content not available locally
- Less incentive to peer with local ISPs



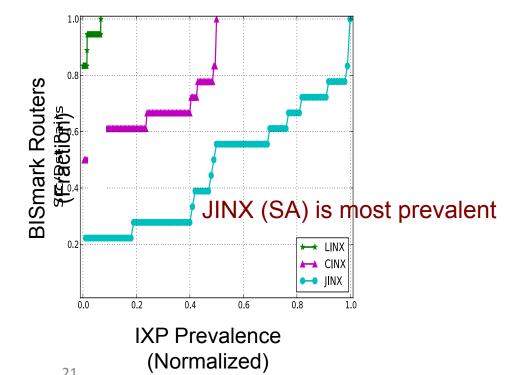
Regional IXPs Only Prevalent on Intra-Country Paths

BISmark !

0.0

Within South Africa:
High Fraction of Paths Have at
Least one Major Regional IXP

Between South Africa and Kenya: Few Paths have Regional IXPs



M-Lab Johannesburg

IXP Prevalence (Normalized) M-Lab Nairobi

LINX (UK) is most prevalent

★ LINX

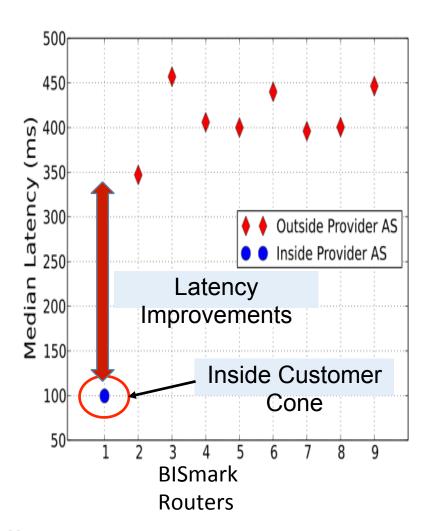
►► AMS-IX

Two Questions

• What is the nature of Internet interconnectivity (between ISPs) in Africa?

 What can be done to reduce latency to common Internet services?

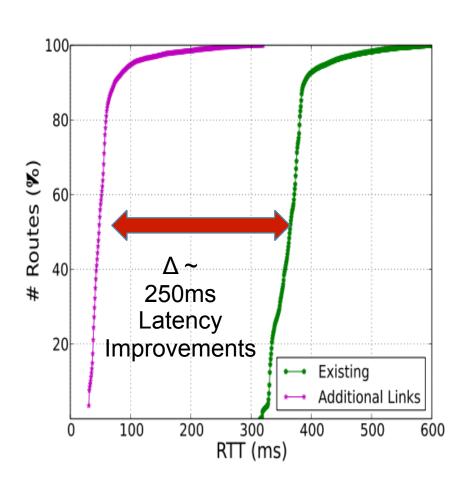
Solution #1: Add More Caches



- Traceroute Probes between BISmark routers (eyeball) and Google Cache Node in Uganda (content)
- Google cache hosted by MTN
- Emulates scenario where content is in nearby country

Latency improvements are limited when peering to the cache is not adequate.

Solution #2: Add More Peering Links

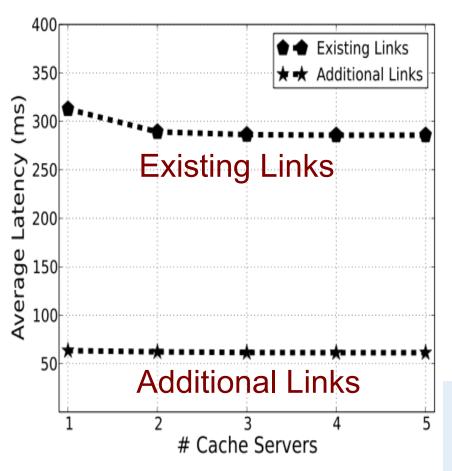


- Simulation: Add peering links between all the participants at
 - JINX (Johannesburg)
 - KIXP (Nairobi)
- Emulates scenario where more ISPs connect and peer at regional IXPs

Additional peering links

Significant latency improvements

Better Peering is a Substitute for Additional Caches



• Experiment:

- add caches in Kenya
- traceroute Probe from SA
- Two scenarios
 - Use existing peering links
 - Add more peering links

Additional caches have little effect on average latency (compared to adding more peering links).

Peering in South Africa: Summary

- What is the nature of Internet interconnectivity (between ISPs) in Africa?
 - Many ISPs are not present in regional IXPs
 - Many ISPs do not interconnect at regional IXPs
- What can be done to reduce latency to common Internet services?
 - Peering at regional IXPs can reduce median intra-continent latencies by 250ms
- Next steps: Better mechanisms for interconnectivity

Internet Routing is Not Flexible Enough

- Routing only on destination IP address blocks
 (No customization of routes by application or sender)
- Can only influence immediate neighbors (No ability to affect path selection remotely)
- Indirect control over packet forwarding (Indirect mechanisms to influence path selection)
- Enables only basic packet forwarding (Difficult to introduce new in-network services)

Valuable Wide-Area Services

- Application-specific peering
 - Route video traffic one way, and non-video another
- Blocking denial-of-service traffic
 - Dropping unwanted traffic further upstream
- Server load balancing
 - Directing client requests to different data centers
- Steering through network functions
 - Transcoders, scrubbers, caches, crypto, ...
- Inbound traffic engineering
 - Splitting incoming traffic over multiple peering links

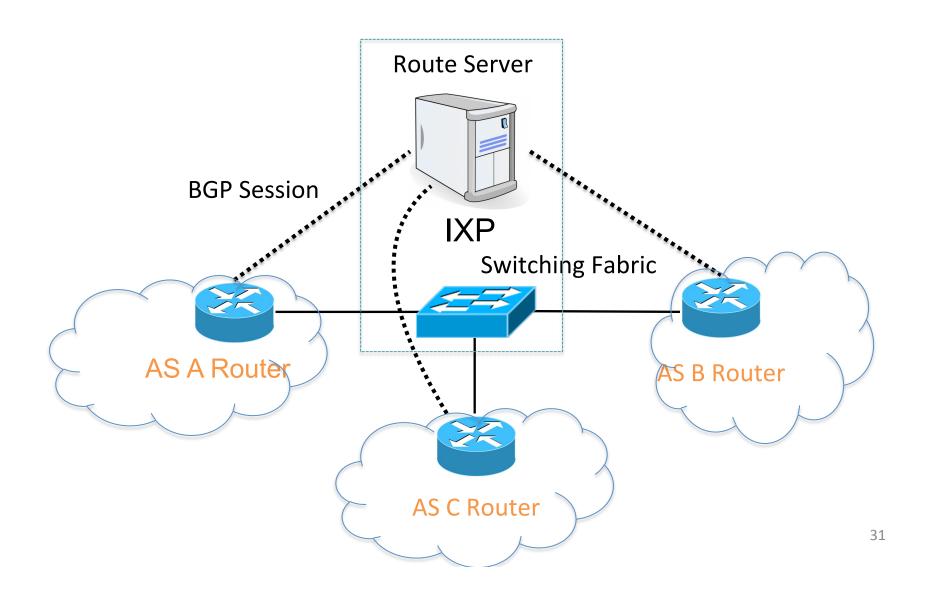
Software Defined Networking (SDN)

- Match packets on multiple header fields (not just destination IP address)
- Control entire networks with a single program (not just immediate neighbors)
- Direct control over packet handling (not indirect control via routing protocol arcana)
- Perform many different actions on packets (beyond basic packet forwarding)

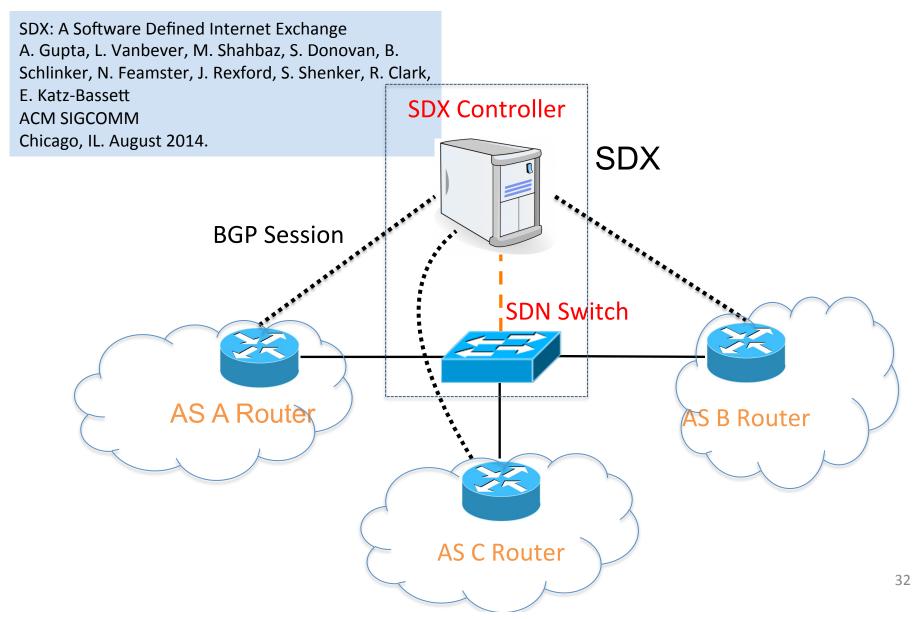
Deploy SDN at Internet Exchanges

- Leverage: SDN deployment even at single IXP can benefit tens to hundreds of providers
 - Without providers deploying new equipment!
- Innovation hotbed: Incentives to innovate, as IXPs on front line of peering disputes
- Growing in numbers:
 - 350-400 IXPs
 - ~100 new IXPs established in past few years

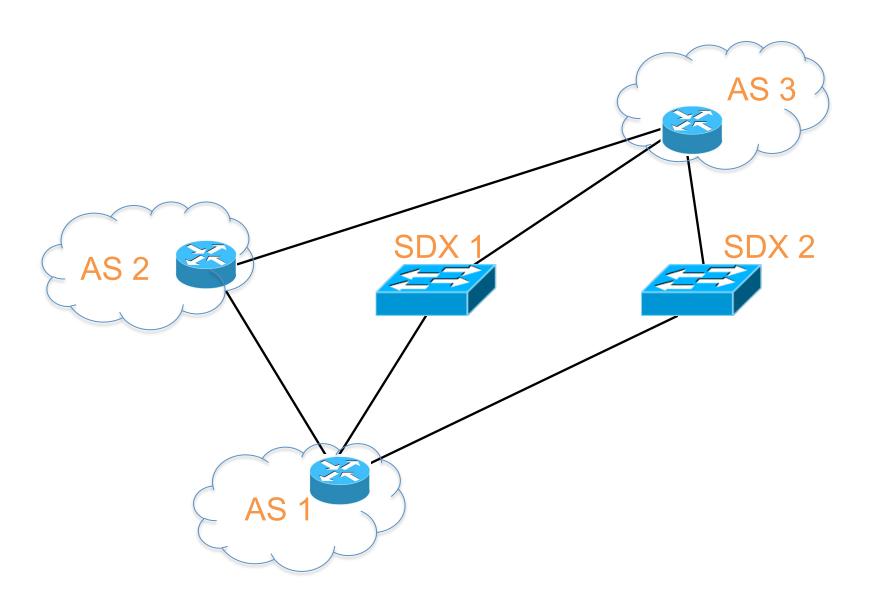
Conventional IXPs



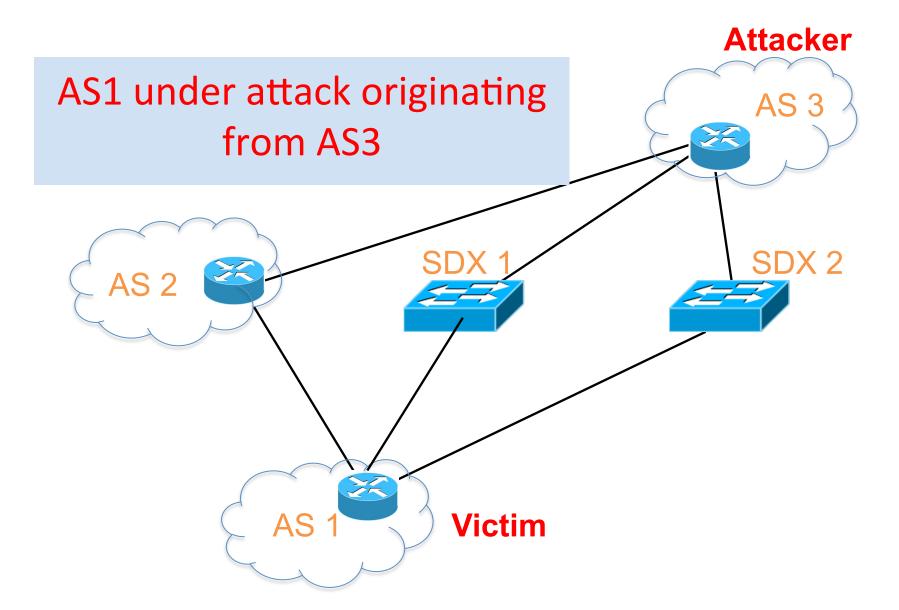
SDX = SDN + IXP



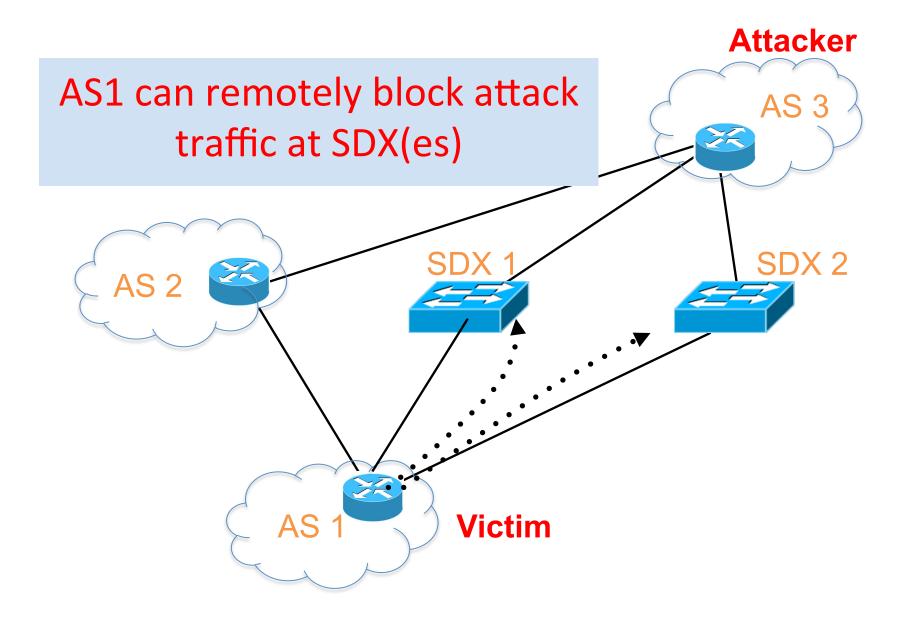
Prevent DDoS Attacks



Prevent DDoS Attacks



Use Case: Prevent DDoS Attacks



SDX-based DDoS protection vs. Traditional Defenses/Blackholing

Remote influence

Physical connectivity to SDX not required

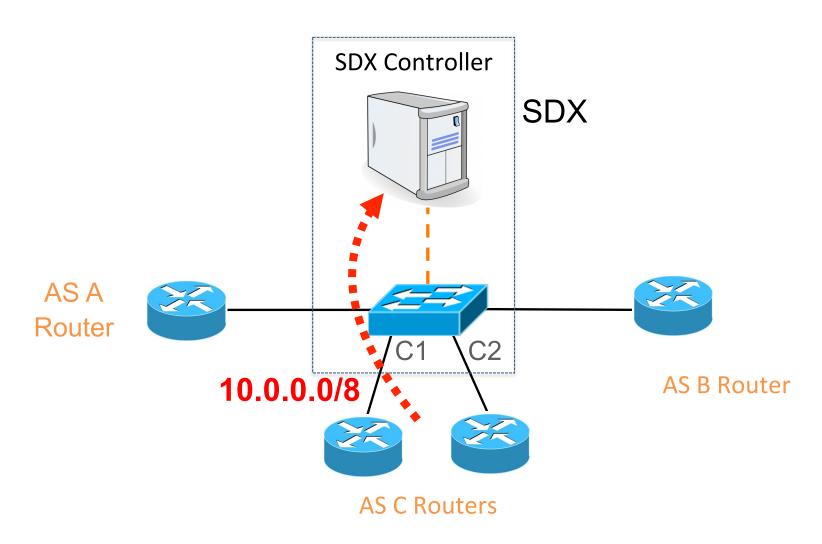
More specific

Drop rules based on multiple header fields, source address, destination address, port number ...

Coordinated

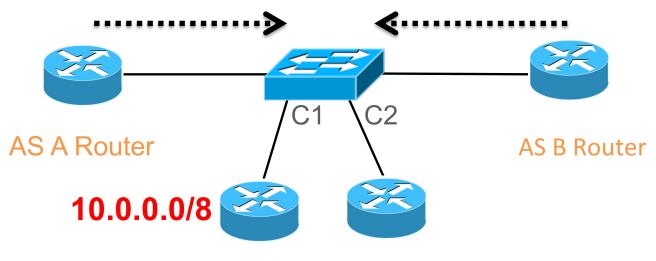
Drop rules can be coordinated across multiple IXPs

Inbound Traffic Control



Inbound Traffic Control

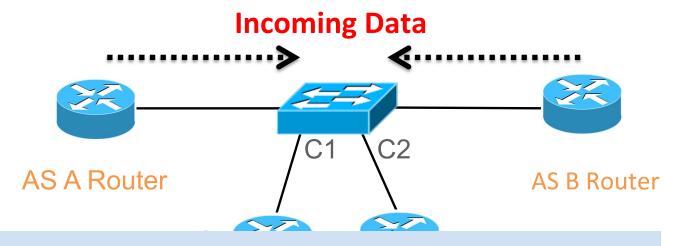
Incoming Data



AS C Routers

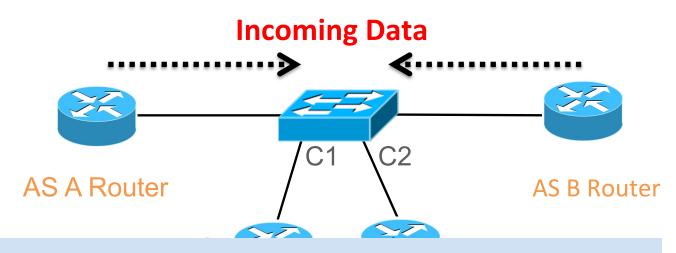
Incoming Traffic	Out Port	Using BGP	Using SDX
dstport = 80	C1		

Inbound Traffic Control



Fine grained policies not possible with BGP

Incoming Traffic	Out Port	Using BGP	Using SDX
dstport = 80	C1	?	



Enables fine-grained traffic engineering policies

Incoming Traffic	Out Port	Using BGP	Using SDX
dstport = 80	C1	?	match(dstport =80)→ fwd(C1)

Building SDX is Challenging

Programming abstractions

How networks define SDX policies and how are they combined together?

Interoperation with BGP

How to provide flexibility w/o breaking global routing?

Scalability

How to handle policies for hundreds of peers, half million address blocks, and matches on multiple header fields?

Building SDX is Challenging

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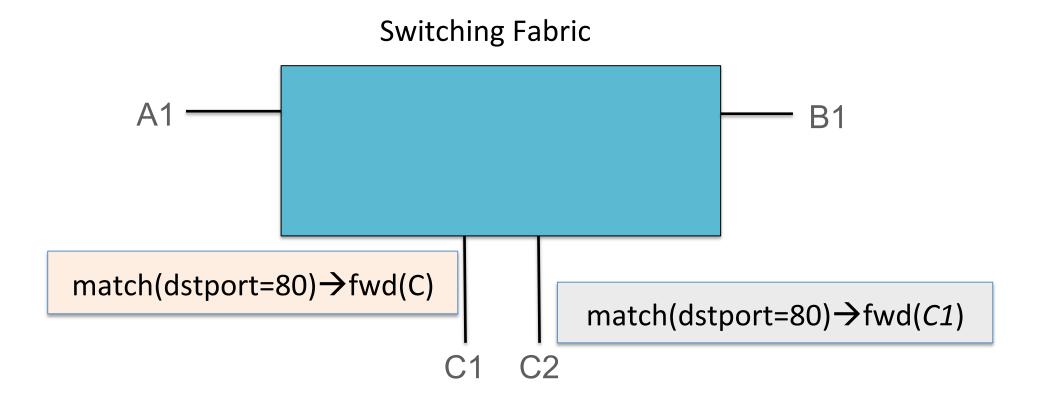
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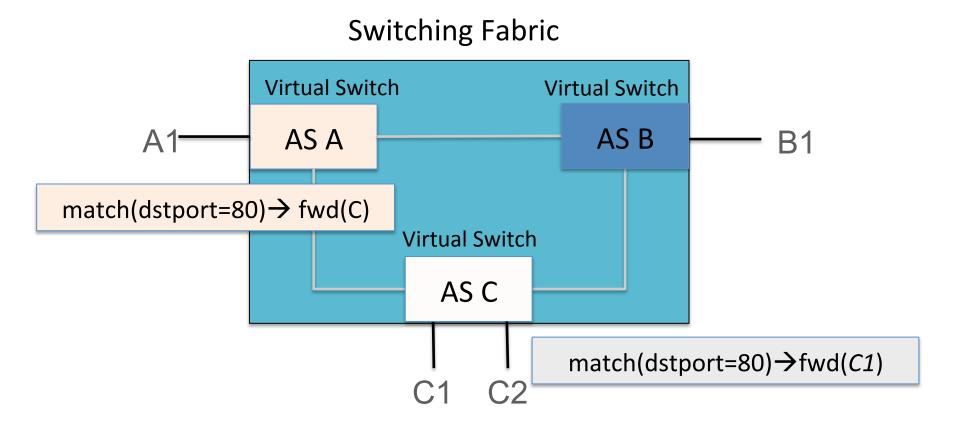
How to handle policies for hundreds of peers, half million prefixes and matches on multiple header fields?

Directly Program the SDX Switch



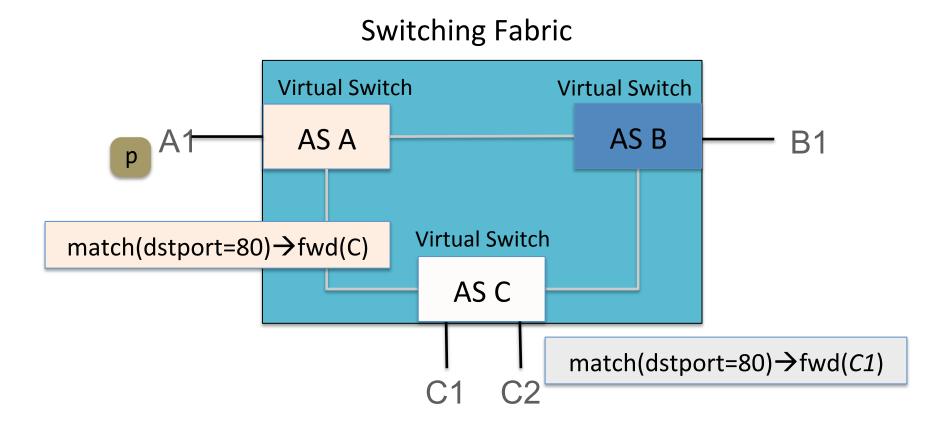
AS A & C directly program the SDX Switch

Virtual Switch Abstraction



Each AS writes policies for its own virtual switch

Combining Participant's Policies



Synthesize: match(inport=A1 & dstport=80) → fwd(C1)

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SDX Platform

- Running code with full BGP integration
 - Github available from http://sdx.cs.princeton.edu
- SDX testbeds:
 - Transit Portal for "in the wild" experiments
 - Mininet for controller experiments
- Exploring deployment opportunities
 - Princeton, DOD/IC, GENI, SOX, Internet2, ESnet
 - Regional IXPs in US, Europe, and Africa

SDX: Summary

- The Internet is changing
 - New challenges for content delivery
 - Increasing importance of IXPs
- SDN can let providers innovate
 - New capabilities and abstractions
- Next steps
 - Operational deployments
 - Additional SDX applications
 - Distributed exchange points

Getting Involved

- UCT Honors Project (To Be Announced)
 - Mobile Measurement
 - Performance, Throttling, Censorship
- Collaborations with UCT Networking Lab

Summer Internship at Princeton

Come for an Internship!

